

Asymptotic notation-

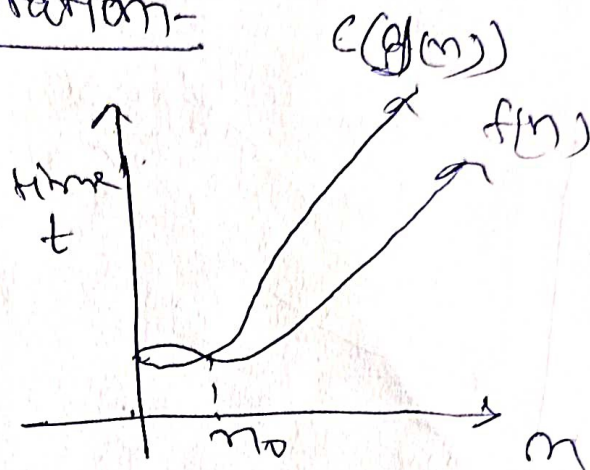
① Big(Oh) - 'O'

$$f(n) \leq c g(n)$$

$$n \geq n_0$$

$$c > 0 \quad n_0 \geq 1$$

$$f(n) = O(g(n))$$



where c, n are real no.

Ex

$$f(n) = 3n + 2$$

$$g(n) = n$$

$$f(n) = O(g(n))$$

$$f(n) \leq c g(n)$$

where

$$c \geq 0$$

$$n_0 \geq 1$$

$$3n + 2 \leq cn$$

Let $c = 4$ $3n + 2 \leq 4n$

$$n \geq 2$$

$g(n)$ will be $n^2, n^3, n^4, 2^n$

but n is suitable and useful

so $f(n) = O(n)$

Remark

$$1 < \log n < \sqrt{n} < n < n \log n < n^2 < n^3 < \dots < 2^n < 3^n < n^n$$

Lower Bound $f(n) = 3n + 2$

Upper Bound

→ closest fun

$$3n + 2 \leq 10n \text{ True}$$

$$\leq 7n \text{ True}$$

$$\leq 100n \text{ "}$$

$$3n + 2 \leq 3n + 2n$$

$$3n + 2 \leq 5n \text{ True}$$

$$\leq 2n^2 + 3n^2 \text{ "}$$

$$\leq 5n^2$$

① "

$$f(n) = O(n) \text{ True}$$

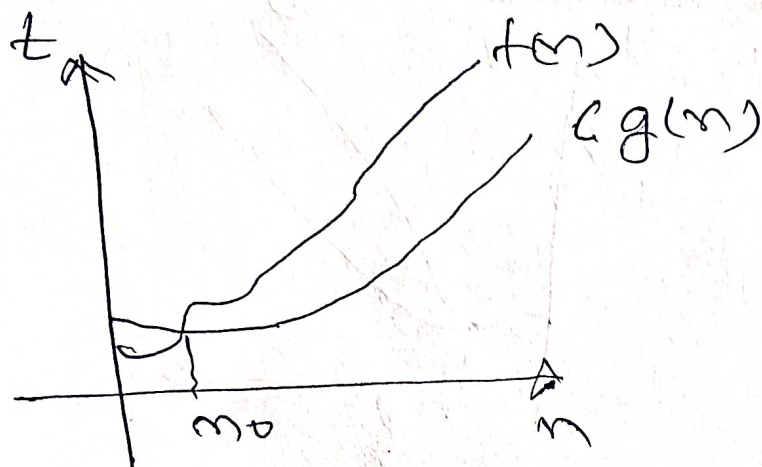
$$f(n) = O(n^2) \text{ True}$$

$$f(n) = O(n^3) \text{ "}$$

$$f(n) = O(2^n) \text{ "}$$

$$f(n) = O(\log n) \text{ False}$$

Big Omega: Ω



$$\boxed{f(n) = \Omega(g(n))}$$

if $f(n) \geq c g(n)$
where $c > 0, n_0 \geq 1$
 $n > n_0$

Ex if $f(n) = 3n + 2$ $g(n) = n$

$$\therefore f(n) \geq c g(n)$$

$$3n + 2 \geq 1 \cdot n \quad \boxed{c=1}$$

Here $3n + 2 = \Omega(n)$

Note -

Let $f(n) = 3n + 2$ $g(n) = n^2$

check that whether it
satisfy Ω condition

Solⁿ $3n + 2 \geq c n^2$

if $c = 1$ it not satisfy

it satisfy only values
less than n like
 $\log n, \log \log n$



$$f(n) = O(g(n))$$

$$\text{if } c_1 g(n) \leq f(n) \leq c_2 g(n)$$

$$c_1, c_2 > 0, n > n_0$$

$$\text{and } n_0 \geq 1$$

Ex $f(n) = 3n + 2$ $g(n) = n$

(a) $f(n) \leq c_2 g(n)$

$$3n + 2 \leq 4n \quad \therefore c_2 = 4$$

(b) $f(n) \geq c_1 g(n)$

$$3n + 2 \geq 1n$$

both condition a & b are satisfy

so $f(n) = O(g(n))$

Ex

$$1 < \log n < \sqrt{n} < n < n \log n < n^2 < \dots < 2^n < 3^n < \dots < n^n$$

(a) $f(n) = 2n^2 + 3n + 4$

$$2n^2 + 3n + 4 \leq 2n^2 + 3n^2 + 4n^2$$



$$f(n)$$

$$f(n) = O(n^2)$$

$$\leq 9n^2 \text{ for } n \geq 1$$



$$c$$



$$g(n)$$

⑤ if $2n^2 + 3n + 4 \geq 1 \times n^2$

$$\begin{array}{ccc} \downarrow & & \downarrow \downarrow \\ f(n) & & c \quad g(n^2) \end{array}$$

$f(n) \geq c g(n^2)$ so twice $\boxed{\Omega(n^2)}$

⑥

$$1 \times n^2 \leq 2n^2 + 3n + 4 \leq 9n^2$$

this is $\boxed{\Theta(n^2)}$

Ex $f(n) = n^2 \log n + n$

$$1 \times n^2 \log n \leq n^2 \log n \leq 10 n^2 \log n$$

so it is $\Theta(n^2 \log n)$ $\Omega(n^2 \log n)$

and $\Theta(n^2 \log n)$

we can classify it between n^2 and n^3

Ex $f(n) = n!$

$$\Rightarrow n(n-1)(n-2) \dots 3 \times 2 \times 1$$

$$1 \times 1 \times 1 \leq 1 \times 2 \times 3 \times \dots \times n \leq n \times n \times n \times \dots \times n$$

$$1 \leq n! \leq n^n$$

it is only $\boxed{\Theta(n^n)}$
 $\Omega(1)$

$\therefore$ both sides of $n!$ are not same so we cannot find Θ

Proof

$$f(n) = \log n!$$

$$\log(1 \times 2 \times \dots \times 1) \leq \log(1 \times 2 \times 3 \times \dots \times n) \leq \log(n \times n \times \dots \times n)$$

$$1 \leq \log n! \leq \log n^n$$

$\Rightarrow$

$$1 \leq \log n! \leq n \log n$$

$$O(n \log n)$$

Time